Principles of Computer Science II Sorting Algorithms

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Lecture 6

Selection Sort Algorithm

This algorithm first finds the smallest element in the array and exchanges it with the element in the first position, then find the second smallest element and exchange it with the element in the second position, and continues in this way until the entire array is sorted.







Selection Sort: Example

Original Array	After 1st pass	After 2nd pass	After 3rd pass	After 4th pass	After 5th pass
3	1	1	1	1	1
6	6	3	3	3	3
0	3 [6	4	4	4
8	8	8	8	5	5
4	4	4	6	6	6
5	5	5	<u>6</u>	8	8

```
1 a = [5, 1, 6, 2, 4, 3]
2 for i in range(0, len(a)):
3     min = i
4     for j in range(i + 1, len(a) - 1):
5         if a[j] > a[min]:
6         min = j
7
8     temp = a[j]
9     a[j] = a[min]
10     a[min] = temp
```









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- ► How much memory is needed ?
 - ▶ 1 additional slot.





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Bubble Sort Algorithm

Bubble Sort is an algorithm which is used to sort N elements that are given in a memory. Bubble Sort compares all the element one by one and sort them based on their values.

- ▶ It is called Bubble sort, because with each iteration the largest element in the list bubbles up towards the last place, just like a water bubble rises up to the water surface.
- ▶ Sorting takes place by stepping through all the data items one-by-one in pairs and comparing adjacent data items and swapping each pair that is out of order.

Bubble Sorting: Example

l	•)	1		•	l
	5					3	
	1	5	6	2	4	3	٦
	1	5	2	6	4	3	7
	1		2		6	3	J

Lets take this Array.

Here we can see the Array after the first iteration.

Similarly, after other consecutive iterations, this array will get sorted.







Bubble Sort Code

► The above algorithm is not efficient because as per the above logic, the for-loop will keep executing for six iterations even if the list gets sorted after the second iteration.







- ▶ We can insert a flag and can keep checking whether swapping of elements is taking place or not in the following iteration.
- ▶ If no swapping is taking place, it means the list is sorted and we can jump out of the for loop, instead executing all the iterations.



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 - N comparisons are required







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Insert Sort Algorithm

larger lists.

increases its efficiency.

additional memory space.

one by one.

 $(n-1)+(n-2)+(n-3)+\ldots+3+2+1$ comparisons are required

A simple Sorting algorithm which sorts the list by shifting elements

▶ It is efficient for smaller data sets, but very inefficient for

▶ Insertion Sort is adaptive, that means it reduces its total

number of steps if given a partially sorted list, hence it

▶ It is better than Selection Sort and Bubble Sort algorithms.

▶ Like Bubble Sorting, insertion sort also requires a single

▶ It has one of the simplest implementation

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- ▶ How much memory is needed ?
 - ▶ 1 additional slot.







Insertion Sort: Example

(Always we start with the second element as key.)

Lets take this Array.

As we can see here, in insertion sort, we pick up a key, and compares it with elemnts ahead of it, and puts the key in the right place

5 has nothing before it.

1 is compared to 5 and is inserted before 5.

6 is greater than 5 and 1.

2 is smaller than 6 and 5. but greater than 1, so its is inserted after 1.

And this goes on...









Insertion Sort Code

```
1 a = [5, 1, 6, 2, 4, 3]
2 for i in range(1, len(a)):
3     key = a[i]
4     j = i - 1
5     while j >= 0 and key < a[j]:
6         a[j+1] = a[j]
7     j -= 1
8     a[j+1] = key</pre>
```

- ▶ key: we put each element of the list, in each pass, starting from the second element: a[1].
- ▶ using the while loop, we iterate, until *j* becomes equal to zero or we find an element which is greater than key, and then we insert the key at that position.





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Quick Sort Algorithm

Quick sort is very fast and requires very less additional space. It is based on the rule of Divide and Conquer. This algorithm divides the list into three main parts:

- ▶ Elements less than the Pivot element
- Pivot element(Central element)
- ▶ Elements greater than the pivot element
- ► Sorts any list very quickly
- ▶ Performance depends on the selection of the Pivot element

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- ▶ Is there a "simple" case ?
 - ▶ The list is already sorted
 - ▶ N comparisons are required
- ▶ How much memory is needed ?
 - ▶ 1 additional slot.





Quick Sort: Example

List: 25 52 37 63 14 17 8 6

- ▶ We pick 25 as the pivot.
- ▶ All the elements smaller to it on its left,
- ▶ All the elements larger than to its right.
- ▶ After the first pass the list looks like:

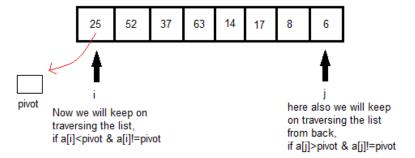
6 8 17 14 25 63 37 52

▶ Now we sort two separate lists:

6 8 17 14 and 63 37 52

▶ We apply the same logic, and we keep doing this until the complete list is sorted.

Quick Sort: Example



if both sides we find the element not satisfying their respective conditions, we swap them. And keep repeating this.

DIVIDE AND CONQUER - QUICK SORT









Quick Sort Code

```
1 a = [25, 52, 37, 63, 14, 17, 8, 6]
3 def partition(list, p, r):
      pivot = list[p]
      i = p
      i = r
      while (1):
           while(list[i] < pivot and list[i] != pivot):
               i += 1
10
          while (list [j] > pivot and list [j] != pivot):
11
              j -= 1
12
13
          if(i < j):
14
               temp = list[i]
15
               list[i] = list[i]
16
               list[j] = temp
17
           else:
18
19
               return j
```

Quick Sort Code

```
1 def quicksort(list, p, r):
2     if (p < r):
3          q = partition(list, p, r)
4          quicksort(list, p, q);
5          quicksort(list, q + 1, r);
6
7 print("Before: ", a)
8 quicksort(a, 0, len(a) - 1)
9 print("After: ", a)</pre>
```







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- ▶ What if we choose the smallest or the largest item as pivot?







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- ▶ What if we choose the median item as pivot?
 - ▶ 1^{st} loop: two lists $\frac{n}{2}$ each
 - ▶ 2^{nd} loop: four lists $\frac{n}{4}$ each
 - **.** . . .
 - ▶ log *n* steps
 - ightharpoonup For each partition we do n comparisons
 - ▶ In total $n \log n$ comparisons







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