Principles of Computer Science II Sorting Algorithms

Ioannis Chatzigiannakis

Sapienza University of Rome

Lecture 8

Selection Sort Algorithm

This algorithm first finds the smallest element in the array and exchanges it with the element in the first position, then find the second smallest element and exchange it with the element in the second position, and continues in this way until the entire array is sorted.



Selection Sort: Example

Original Array	After 1st pass	After 2nd pass	After 3rd pass	After 4th pass	After 5th pass
3	1	1	1	1	1
6	6	3	3	3	3
0	3	6	4	4	4
8	8	8	8	5	5
4	4	4	6	6	6
5	5	5	6	8	8

Selection Sort Code

1 a = [5, 1, 6, 2, 4, 3]2 for i in range(0, len(a)): min = i3 for j in range (i + 1, len(a) - 1): 4 if a[j] > a[min]: 5 6 min = j7 temp = a[j]8 a[j] = a[min]9 10 a[min] = temp



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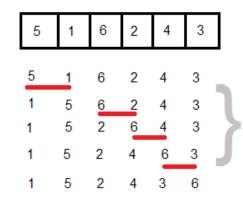
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 - 1 additional slot.

Bubble Sort Algorithm

Bubble Sort is an algorithm which is used to sort N elements that are given in a memory. Bubble Sort compares all the element one by one and sort them based on their values.

- It is called Bubble sort, because with each iteration the largest element in the list bubbles up towards the last place, just like a water bubble rises up to the water surface.
- Sorting takes place by stepping through all the data items one-by-one in pairs and comparing adjacent data items and swapping each pair that is out of order.

Bubble Sorting: Example



Lets take this Array.

Here we can see the Array after the first iteration.

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Similarly, after other consecutive iterations, this array will get sorted.





Bubble Sort Code

```
1a = [5, 1, 6, 2, 4, 3]
2 for i in range(0, len(a)):

3 for j in range(0, len(a) - i - 1):

4 if a[j] > a[j+1]:

5 temp = a[j]

6 a[j] = a[j+1]

7 a[j+1] = temp
```

The above algorithm is not efficient because as per the above logic, the for-loop will keep executing for six iterations even if the list gets sorted after the second iteration.

Bubble Sort Code: Version 2

- We can insert a flag and can keep checking whether swapping of elements is taking place or not in the following iteration.
- If no swapping is taking place, it means the list is sorted and we can jump out of the for loop, instead executing all the iterations.

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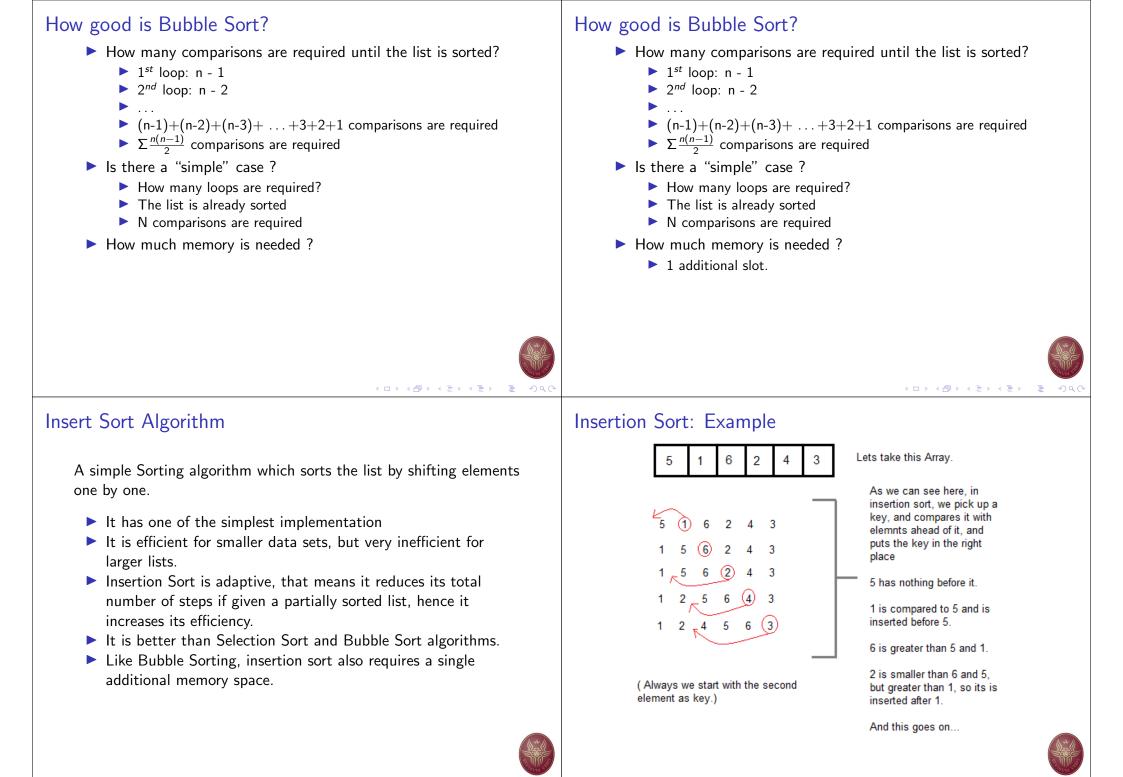
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Insertion Sort Code

```
1 \mathbf{a} = [5, 1, 6, 2, 4, 3]
2 for i in range(1, len(a)):
     key = a[i]
4
     i = i - 1
     while j \ge 0 and key < a[j]:
5
          a[j+1] = a[j]
6
         i -= 1
7
     a[i+1] = key
8
```

- **key**: we put each element of the list, in each pass, starting from the second element: a[1].
- using the while loop, we iterate, until j becomes equal to zero or we find an element which is greater than key, and then we insert the key at that position.

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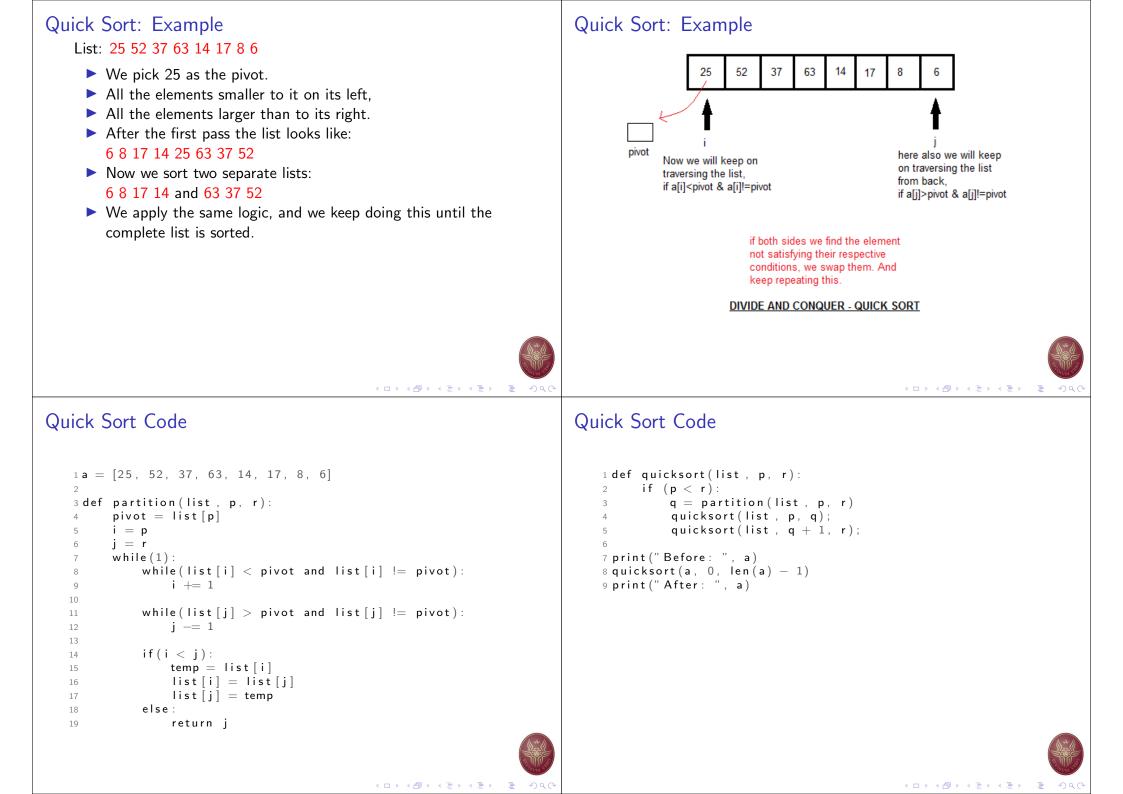
Quick Sort Algorithm

Quick sort is very fast and requires very less additional space. It is based on the rule of Divide and Conquer. This algorithm divides the list into three main parts :

- Elements less than the Pivot element
- Pivot element(Central element)
- Elements greater than the pivot element
- Sorts any list very quickly
- Performance depends on the selection of the Pivot element



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 - ▶ 2^{nd} loop: four lists $\frac{n}{4}$ each
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 - For each partition we do *n* comparisons
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