

# Principles of Computer Science II

## Recursive Algorithms

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### Lecture 9



# Recursion Coding Style

Recursion is a way of programming or coding a problem, in which a function calls itself one or more times in its body. Usually, it is returning the return value of this function call. If a function definition fulfils the condition of recursion, we call this function a recursive function.

Termination condition:

- ▶ A recursive function has to terminate to be used in a program.
- ▶ A recursive function terminates, if with every recursive call the solution of the problem is downsized and moves towards a base case.
- ▶ A base case is a case, where the problem can be solved without further recursion.



# Factorial Computation: Using Iteration

```
1 def iterative_factorial(n):  
2     result = 1  
3     for i in range(2, n+1):  
4         result *= i  
5     return result
```



# Factorial Computation: Using Recursion

```
1 def factorial(n):  
2     if n == 1:  
3         return 1  
4     else:  
5         return n * factorial(n-1)
```



## Factorial Computation

```
1 def factorial(n):
2     print("factorial has been called with n = " + str(n))
3     if n == 1:
4         return 1
5     else:
6         res = n * factorial(n-1)
7         print("intermediate result for ", n, " * factorial
8             (" ,n-1, "): ",res)
9         return res
10 print(factorial(5))
```



## Fibonacci Numbers

The Fibonacci numbers are defined by:

$$F_n = F_{n-1} + F_{n-2}$$

where  $F_0 = 0$  and  $F_1 = 1$

▶ 0,1,1,2,3,5,8,13,21,34,55,89, ...



## Factorial Computation: Using Recursion

```
1 def fib(n):
2     if n == 0:
3         return 0
4     elif n == 1:
5         return 1
6     else:
7         return fib(n-1) + fib(n-2)
```



## Factorial Computation: Using Iteration

```
1 def fibi(n):
2     a, b = 0, 1
3     for i in range(n):
4         a, b = b, a + b
5     return a
```

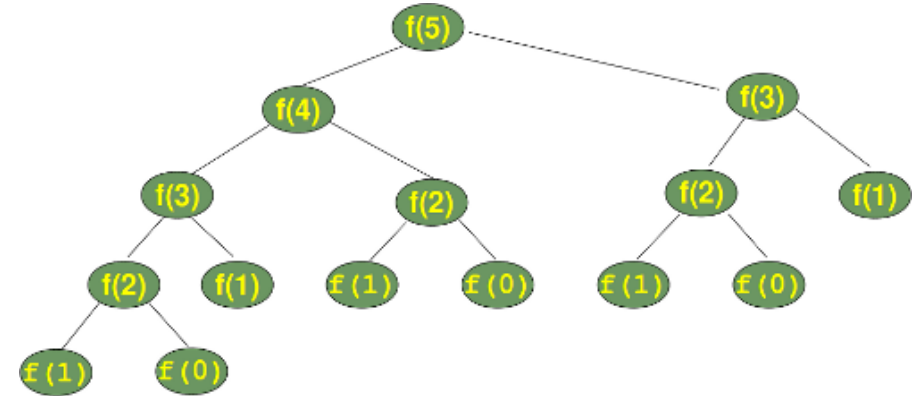


## Measure Performance

```
1 from timeit import Timer
2 from fibo import fib
3
4 t1 = Timer(" fib(10)", "from fibo import fib")
5
6 for i in range(1,41):
7     s = " fib(" + str(i) + ")"
8     t1 = Timer(s, "from fibo import fib")
9     time1 = t1.timeit(3)
10    s = " fibi(" + str(i) + ")"
11    t2 = Timer(s, "from fibo import fibi")
12    time2 = t2.timeit(3)
13    print("n=%2d, fib: %8.6f, fibi: %7.6f, percent: %10.2f"
14          " % (i, time1, time2, time1/time2))
```



## Fibonacci Numbers



## Factorial Computation: Using Recursion and Memory

```
1 memo = {0:0, 1:1}
2 def fibm(n):
3     if not n in memo:
4         memo[n] = fibm(n-1) + fibm(n-2)
5     return memo[n]
```



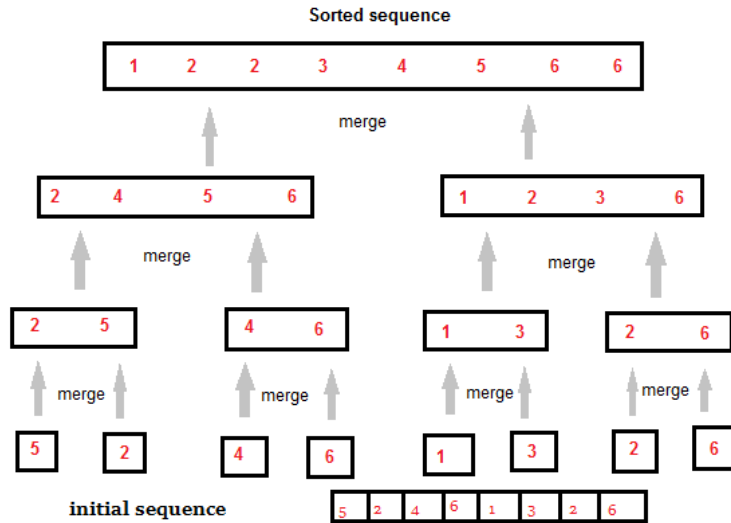
## Merge Sort Algorithm

In Merge Sort the unsorted list is divided into  $N$  sublists, each having one element, because a list consisting of one element is always sorted. Then, it repeatedly merges these sublists, to produce new sorted sublists, and in the end, only one sorted list is produced.

- ▶ Divide and Conquer algorithm
- ▶ Performance always same for Worst, Average, Best case



## Merge Sort: Example



## Merge Sort Code

```

1 a = [25, 52, 37, 63, 14, 17, 8, 6]
2
3 def mergesort(list):
4     if len(list) == 1:
5         return list
6
7     left = list[0: len(list) // 2]
8     right = list[len(list) // 2:]
9
10    left = mergesort(left)
11    right = mergesort(right)
12
13    return merge(left, right)

```



## Merge Sort Code

```

1 def merge(left, right):
2     result = []
3     while len(left) > 0 and len(right) > 0:
4         if left[0] <= right[0]:
5             result.append(left.pop(0))
6         else:
7             result.append(right.pop(0))
8
9     while len(left) > 0:
10        result.append(left.pop(0))
11
12    while len(right) > 0:
13        result.append(right.pop(0))
14
15    return result
16
17 print(" Before: ", a)
18 r = mergesort(a)
19 print(" After: ", r)

```



## How good is Merge Sort?

- ▶ How many comparisons are required until the list is sorted?
  - ▶ 1<sup>st</sup> loop: two lists  $\frac{n}{2}$  each
  - ▶ 2<sup>nd</sup> loop: four lists  $\frac{n}{4}$  each
  - ▶ ...
  - ▶ log  $n$  steps
  - ▶ For each partition we do  $n$  comparisons
  - ▶ In total  $n \log n$  comparisons
- ▶ How much memory is needed?
  - ▶ 1 additional slot.



## Quick Sort Algorithm

Quick sort is very fast and requires very less additional space. It is based on the rule of **Divide and Conquer**. This algorithm divides the list into three main parts :

- ▶ Elements less than the Pivot element
- ▶ Pivot element(Central element)
- ▶ Elements greater than the pivot element
  
- ▶ Sorts any list very quickly
- ▶ Performance depends on the selection of the Pivot element



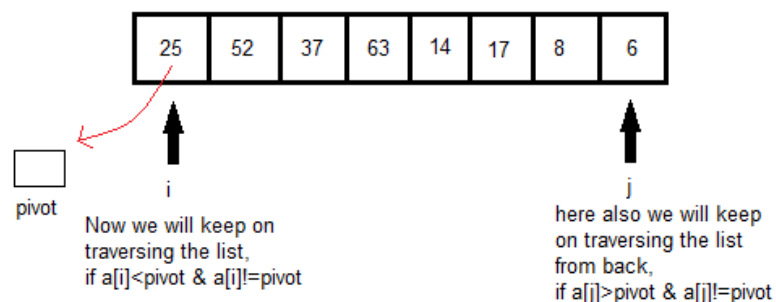
## Quick Sort: Example

List: 25 52 37 63 14 17 8 6

- ▶ We pick 25 as the pivot.
- ▶ All the elements smaller to it on its left,
- ▶ All the elements larger than to its right.
- ▶ After the first pass the list looks like:  
6 8 17 14 25 63 37 52
- ▶ Now we sort two separate lists:  
6 8 17 14 and 63 37 52
- ▶ We apply the same logic, and we keep doing this until the complete list is sorted.



## Quick Sort: Example



if both sides we find the element not satisfying their respective conditions, we swap them. And keep repeating this.

**DIVIDE AND CONQUER - QUICK SORT**



## Quick Sort Code

```
1 a = [25, 52, 37, 63, 14, 17, 8, 6]
2
3 def partition(list, p, r):
4     pivot = list[p]
5     i = p
6     j = r
7     while(1):
8         while(list[i] < pivot and list[i] != pivot):
9             i += 1
10
11        while(list[j] > pivot and list[j] != pivot):
12            j -= 1
13
14        if(i < j):
15            temp = list[i]
16            list[i] = list[j]
17            list[j] = temp
18        else:
19            return j
```



## Quick Sort Code

```
1 def quicksort(list, p, r):
2     if (p < r):
3         q = partition(list, p, r)
4         quicksort(list, p, q);
5         quicksort(list, q + 1, r);
6
7 print(" Before: ", a)
8 quicksort(a, 0, len(a) - 1)
9 print(" After: ", a)
```



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  - ▶ 1<sup>st</sup> loop:  $n - 1$
  - ▶ 2<sup>nd</sup> loop:  $n - 2$
  - ▶ ...
  - ▶  $(n-1) + (n-2) + (n-3) + \dots + 3 + 2 + 1$  comparisons are required
  - ▶  $\sum \frac{n(n-1)}{2}$  comparisons are required



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- ▶ What if we choose the median item as pivot?



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