## Principles of Computer Science II

Sequence Similarity

Ioannis Chatzigiannakis

Sapienza University of Rome

## Lecture 19

## Sequence Similarity

Alignment of ATATATAT vs TATAAT

$$
\begin{array}{cccccccc}
\mathrm{A} & \mathrm{~T} & \mathrm{~A} & \mathrm{~T} & \mathrm{~A} & \mathrm{~T} & \mathrm{~A} & \mathrm{~T} \\
& : & : & : & : & : & : & : \\
- & \mathrm{T} & \mathrm{~A} & \mathrm{~T} & \mathrm{~A} & - & \mathrm{A} & \mathrm{~T}
\end{array}
$$

## Edit Distance

- We looked for repeating patterns within DNA sequences.
- How can we measure the similarity between different sequences?
- We use the notion of Vladimir Levenshtein introduced in 1966
- Edit distance - the minimum number of editing operations needed to transform one string into another (insert/delete symbol or substitute one symbol for another).

Alignment of ATATATAT vs TATATATA
A T A T A T A T -

- T A T A T A T A


## Sequence Similarity

| Alignment of TGCATAT vs ATCCGAT |  |
| :---: | :--- |
| TGCATAT |  |
| $\downarrow$ | delete last T |
| TGCATA |  |
| $\downarrow$ | delete last A |
| TGCAT |  |
| $\downarrow$ | insert A at the front |
| ATGCAT |  |
| $\downarrow$ | substitute C for G in the third position |
| ATCCAT |  |
| $\downarrow$ | insert a G before the last A |
| ATCCGAT |  |

Five operations.

## Sequence Similarity

```
Alignment of TGCATAT vs ATCCGAT
    TGCATAT
            insert A at the front
ATGCATAT
    d delete T in the sixth position
    ATGCAAT
        \downarrow}\quad\mathrm{ substitute G for A in the fifth position
    ATGCGAT
        \downarrow substitute C for G in the third position
ATCCGAT
```

Four operations.

## Edit Distance Algorithm using Dynamic Programming

- Assume two strings:
- $v$ (of $n$ characters)
- $w$ (of $m$ characters)
- The alignment of $v, w$ is a two-row matrix such that - first row: contains the characters of $v$ (in order)
- second row: contains the characters of $w$ (in order)
- spaces are interpersed throughout the table.
- Characters in each string appear in order, though not necessarily adjacently.

| A | T | - | G | T | T | A | T | - |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | T | C | G | T | - | A | - | C |

- No column contains spaces in both rows.
- At most $n+m$ columns.


## Edit Distance

- Vladimir Levenshtein defined the notion of Edit distance
- Did not provide an algorithm to compute it.

Edit Distance Algorithm using Dynamic Programming

| A | T | - | G | T | T | A | T | - |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | T | C | G | T | - | A | - | C |

- Matches - columns with the same letter,
- Mismatches - columns with different letters.
- Columns containing one space are called indels
- Space on top row: insertions
- Space on bottom row: deletions
\# matches + \# mismatches + \#indels $<n+m$


## Representing the rows

| $\mathbf{v}$ | A | T | - | G | T | T | A | T | - |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{w}$ | A | T | C | G | T | - | A | - | C |

- One way to represent $v$
- AT-CGTAT-
- One way to represent $w$
- ATCGT-A-C
- Another way to represent $v$
- AT-CGTAT-
- 122345677
- number of symbols of $v$ present up to a given position
- Similarly, to represent $w$
- ATCGT-A-C
- 123455667


## Edit distance graph

- Edit graph: a grid of $n, m$ size.
- The edit graph will help us in calculating the edit distance.
- Alignment: a path from $(0,0)$ to $(n, m)$.
- Every alignment corresponds to a path in the edit graph.
- Diagonal movement at point $i, j$ correspond to column $\binom{v_{i}}{w_{j}}$
- Horizontal movement correspond to column $\binom{-}{w_{j}}$
- Vertical movement correspond to column $\binom{v_{i}}{-}$

Representing the rows

| $\mathbf{v}$ | A | T | - | G | T | T | A | T | - |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{w}$ | A | T | C | G | T | - | A | - | C |


| $\mathbf{v}$ | 1 | 2 | 2 | 3 | 4 | 5 | 6 | 7 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{w}$ | 1 | 2 | 3 | 4 | 5 | 5 | 6 | 6 | 7 |

can be viewed as a coordinate in 2-dimensional $n \times m$ grid:

$$
\binom{0}{0}\binom{1}{1}\binom{2}{2}\binom{2}{3}\binom{3}{4}\binom{4}{5}\binom{5}{5}\binom{6}{6}\binom{7}{6}\binom{7}{7}
$$

The entire alignment is simply a path:
$(0,0) \rightarrow(1,1) \rightarrow(2,2) \rightarrow(2,3) \rightarrow(3,4) \rightarrow(4,5) \rightarrow(5,5) \rightarrow$ $(6,6) \rightarrow(7,6) \rightarrow(7,7)$

Edit distance graph


```
\
```

Profile most-frequent k-mer

```
1def edit_distance(s1, s2)
    m=len(s1)+1
    n=\operatorname{len}(s2)+1
    tbl}={
    for i in range(m): tbl[i,0]=i
    for j in range(n): tbl[0,j]=j
    for i in range(1,m)
        for j in range(1, n):
            cost = 0 if s1[i-1]=s2[j-1] else 1
            tbl[i,j] = min(tbl[i, j - 1]+1, tbl[i-1, j]+1
            tbl[i-1, j -1]+cost)
    return tbl[i,j]
```


## Profile most-frequent k-mer

```
1 def levenshteinDistance(s1, s2):
    if len(s1) > len(s2)
        s1, s2 = s2, s1
    distances = range(len(s1) + 1)
    for i2, c2 in enumerate(s2)
        distances_ = [i2 +1]
        for i1, c1 in enumerate(s1):
            if c1 = c2:
            distances_.append(distances[i1])
            else:
            distances_.append(1 + min((distances [i1]
                distances[i1 + 1], distances_[-1])))
            distances = distances_
    return distances[-1]
```

