Principles of Computer Science II Algorithms for BioInformatics

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Lecture 5





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- Game played in turns by 2 players.
- Two piles of equal number of pebbles.
- ► Each turn a player may either
 - take 1 pebble from a single pile, or
- take 1 pebble from both piles. The player that takes the last pebble wins.



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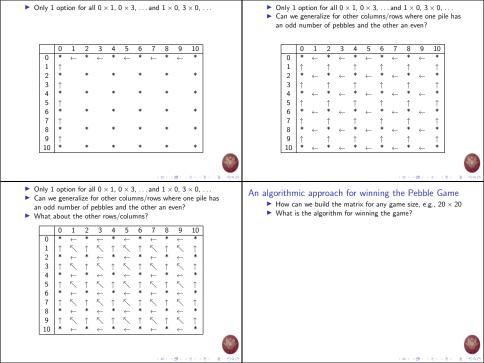
- Does the first player always have an advantage? Let's consider the most simplified version.
 - Pebbles = 2 we call this the 2 x 2 game.

 - Is there a winning strategy? ► What is the winning strategy?

Generaled Strategy for Winning the Pebble Game

- Can we generalize the strategy of the 2 x 2 game?
- ▶ What about the 3 × 3 game? Consider different game sequences.
- Consider the n x n game.
 - Is there only one winning strategy?
 - How easy it is to describe our strategy? Quality of solution.

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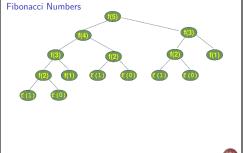


An algorithmic approach for winning the Pebble Game How can we build the matrix for any game size, e.g., 20 × 20 What is the algorithm for winning the game? Why should I care?	An algorithmic approach for winning the Pebble Game How can we build the matrix for any game size, e.g., 20 × 20 What is the algorithm for winning the game? Why should I care? It is the sequence alignment problem. The computational idea used to solve both problems is the same. We need to understand how algorithms work.
Methodology of solving a computational problem What is the problem at hand? Identify & Understand assumptions. What is our goal? Identify similar problems/solutions in the bibliography What are the theoretical foundation? Can we formulate the problem in a unambiguous and precise way? What is the Input that we have? Do we have enough data or should we try to collect? Open data sets? Can we synthesize input data? What is the expected Output?	Solution Sketch Do we have a rough idea of a solution? Do we have identified an approach to solving the problem? think again! po through the definition – maybe we overlooked something? Write down a solution sketch can you try it out with a small input? Is the solution correct? can we provide some arguments? What is the performance of the solution? Can we think of a more efficient solution?

Implement the first version Iterative approach Pick your programming language of choice. Step-by-step development Continuous development. Implement your solution Agile methodology. No need to try to make it elegant / fast. Remember Donalt Knuth: There is no such thing as early Identify issues in previous version optimization. Code beautification. Bug fixes. Get some input data Open datasets Performance improvements. Additional functionalities. Small size ▶ Limited Evaluation Implement improvements does it work ? Make sure code is always clean + easy to maintain. Keep detailed records of changes. do you need to make any modifications ? Always keep history of source code evolution. are there special cases that you missed ? ▶ Performance Evaluation bigger input. scalability ? 40 × 40 × 42 × 42 × 2 × 990 10110 10111 121 121 2 000 Theoretical – Practical Approach Cycle Quality of Code Theoretical Results Modeling Analysis John Woods Always code as if the guy who ends up maintaining your code will be a violent psychopath who knows Algorithms Deployment where you live. Experimental Algorithm Evaluation Engineering Simulation Implementation 4 D > 4 B > 4 B > 4 B > 3 B + 4 9 0 101 (8) (2) (2) (3) 3 (9)

Recursion Coding Style	Factorial Computation: Using Iteration
Recursion is a way of programming or coding a problem, in which a function calls itself one or more times in its body. Usually, it is returning the return value of this function call. If a function definition fulfils the condition of recursion, we call this function a recursive function.	<pre>def iterative_factorial(n): result = 1 for i in range(2,n+i): result == i return result</pre>
Termination condition: ➤ A recursive function has to terminate to be used in a program. ➤ A recursive function terminates, if with every recursive call the solution of the problem is downsized and moves towards a base case. ➤ A base case is a case, where the problem can be solved without further recursion.	
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Factorial Computation: Using Recursion	Factorial Computation
<pre>def factorial(n): if n == 1: return 1 else: return n * factorial(n-1)</pre>	<pre>def factorial(n): print("factorial has been called with n = " + str(n)) if n == 1: return 1 else: res = n * factorial(n-1) print("intermediate result for ", n, " * factorial(" ,n-1, "): return res</pre>
	<pre>print(factorial(5))</pre>
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Fibonacci Numbers	Factorial Computation: Using Recursion
The Fibonacci numbers are defined by: $F_n = F_{n-1} + F_{n-2}$ where $F_0 = 0$ and $F_1 = 1$ \blacktriangleright 0.1.1.2.3,5.8.13,21,34,55,89,	<pre>def fib(n): if n == 0: return 0 elif n == 1: return 1 else: return fib(n-1) + fib(n-2)</pre>
Factorial Computation: Using Iteration	Measure Performance
<pre>def fibi(n): a, b = 0, 1 for i in range(n): a, b = b, a + b return a</pre>	<pre>import time for i in range(1,41): t1 = time.perf_counter() s = fib(i) t2 = time.perf_counter() - t1 t3 = time.perf_counter() s = fib(i) t4 = time.perf_counter() - t3 print(f"n={i}, fib: {t2:.2f}, fibi: {t1:.2f}, percent: {t2/t4:.</pre>
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Factorial Computation: Using Recursion and Memory

```
memo = {0:0, 1:1}
def fibm(n):
    if not n in memo:
        memo[n] = fibm(n-1) + fibm(n-2)
    return memo[n]
```





