

# Pervasive Systems

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Lecture 7:  
DTN Bundle Protocol



# Delay and Disruption Tolerant Networking (DTN)

- DTN overcome problems associated with
  - 1 Intermittent connectivity,
  - 2 Long or variable delays,
  - 3 Assymmetric data rates,
  - 4 High error rates.
- DTN uses the Store and Forward Message Switching.
- Resembles mechanisms invented in the ancient times: Pony-Express, Postal systems.



# Store and Forward Message Switching



- Individual messages are grouped in Bundles.
- Bundles (of messages) are moved from a storage place of one node, to a storage place of another node (switch interaction).
- Store-and-forwarding methods are also used in today's voicemail and email systems.
  - these systems are not node-to-node relays (as shown above) but rather star relays.
  - both the source and destination independently contact a central storage device at the center of the links.



# The need for Persistent Storage

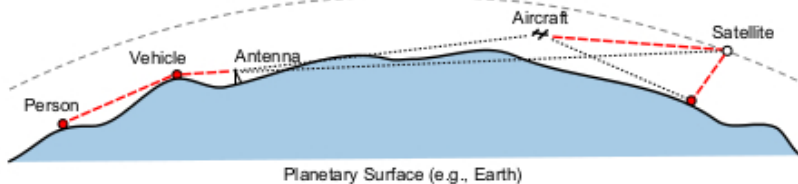
DTN routers need persistent storage for their queues for one or more of the following reasons:

- 1 A communication link to the next hop may not be available for a long time.
- 2 One node in a communicating pair may send or receive data much faster or more reliably than the other node.
- 3 A message, once transmitted, may need to be retransmitted if an error occurs at an upstream (toward the destination) node, or if an upstream node declines acceptance of a forwarded message.



## Opportunistic Contacts

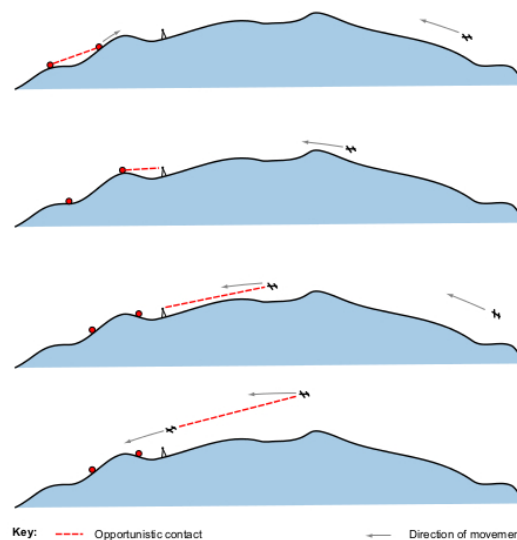
- Network nodes may need to communicate during **opportunistic** contacts.
- Sender and Receiver make contact at an **unscheduled** time.



Key:   
- - - Connected link ● Mobile or fixed wireless surface node ○ Orbiting node   
 ····· Disconnected (obstructed or silent) link



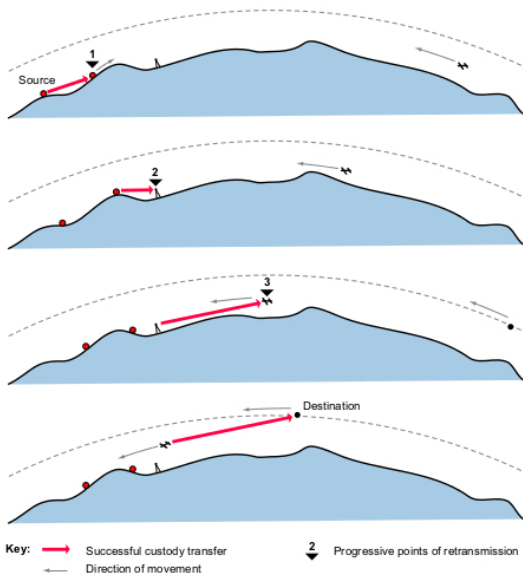
## Opportunistic Contacts: An Example



Key:   
- - - Opportunistic contact   
 ← Direction of movement



## Opportunistic Contacts: An Example

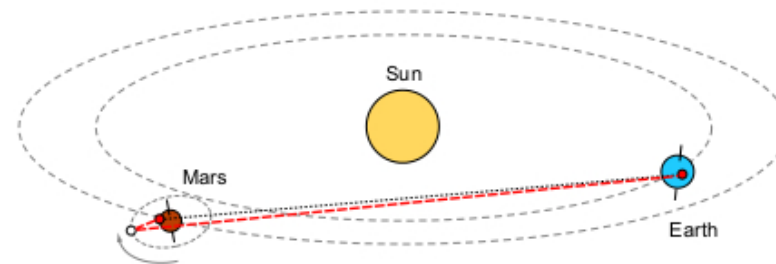


Key:   
→ Successful custody transfer   
 ····· Progressive points of retransmission   
 ← Direction of movement



## Scheduled Contacts

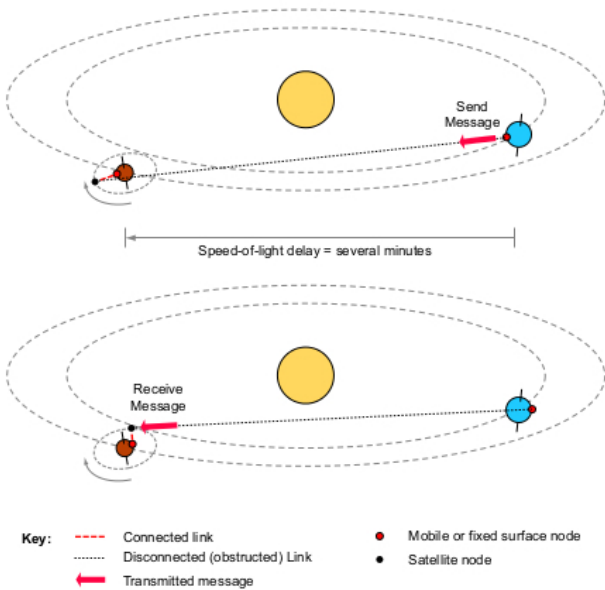
- In many scenarios we can predict (or receive) times schedules of nodes future positions.
- We can arrange future communications sessions.
- Time synchronization between nodes is of crucial importance.



Key:   
- - - Connected link ● Mobile or fixed wireless surface node ○ Orbiting node   
 ····· Disconnected (obstructed or silent) link

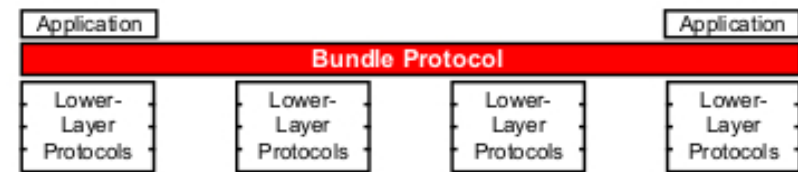


## Scheduled Contacts: An Example

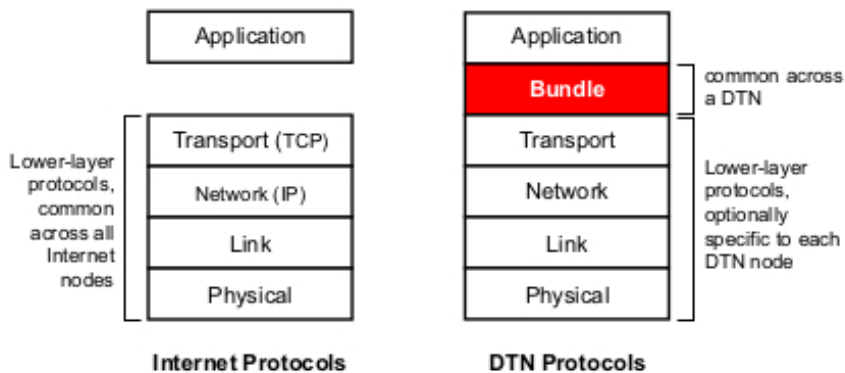


## Main Concepts of Bundle Protocol

- Implements store-and-forward message switching.
- Overlays a new transmission protocol (the bundle protocol) on top of the lower layers (e.g., the Internet protocols).
- Ties together the lower layers so that application programs can communicate across the same or different sets of lower-layer protocols under conditions that involve long network delays or disruptions



## IP stack vs DTN stack

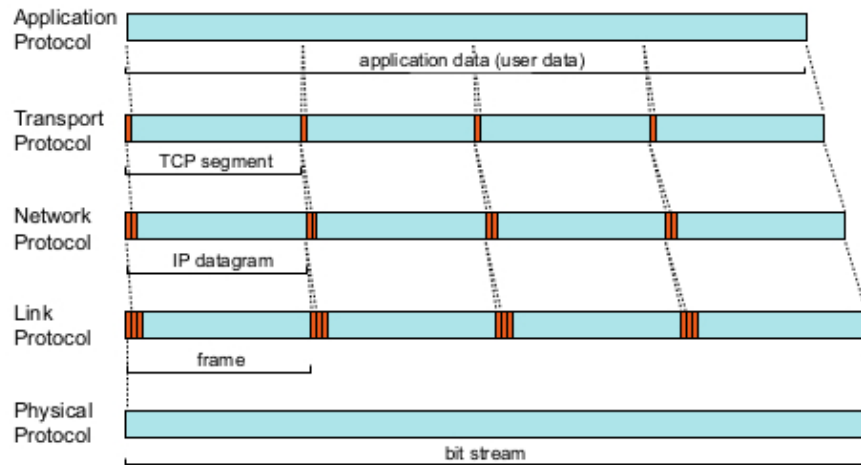


## Bundles

Bundles consist of three things:

- 1 a bundle header consisting of one or more DTN blocks inserted by the bundle-protocol agent,
- 2 a source-applications user data, including control information provided by the source application for the destination application that describes how to process, store, dispose of, and otherwise handle the user data, and
- 3 an optional bundle trailer, consisting of zero or more DTN blocks, inserted by the bundle-protocol agent (not shown in the figure below). Like application-program user data, bundles can be arbitrarily long.

# Bundles Encapsulation



Key:  Header  
 User data

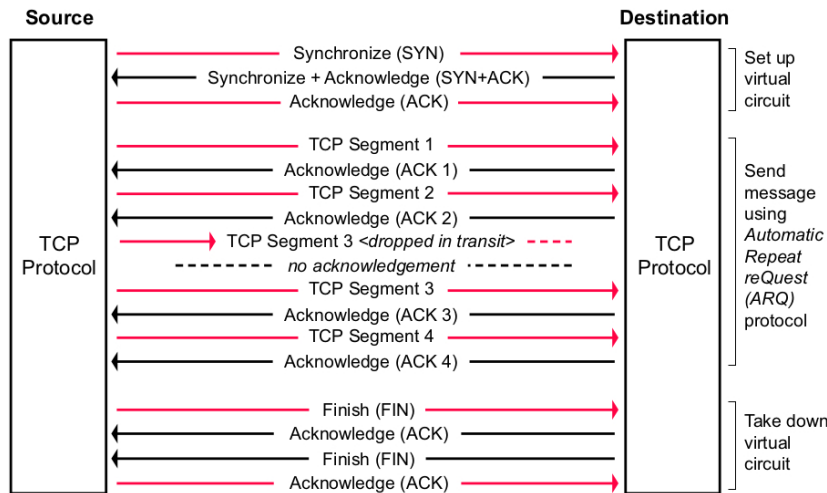


# Today's Internet – Conversational Protocols

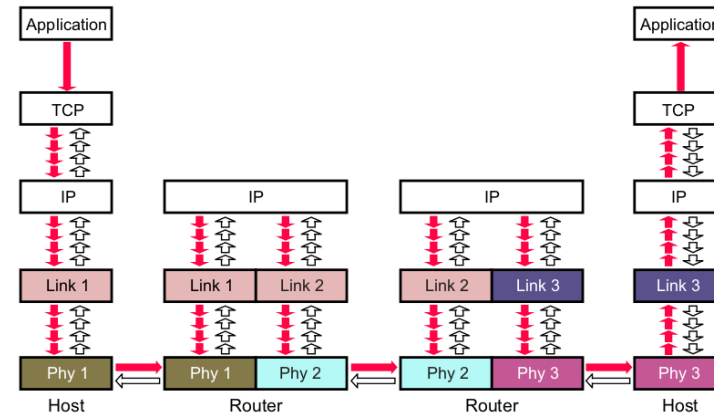
- TCP protocol is said to be conversational (interactive).
- A complete one-way message involves many source-to-destination signaling round-trips:
  - 1 **Set Up:** A three-way “Hello” handshake.
  - 2 **Segment Transfer and Acknowledgement:** Each TCP segment (or a few segments) sent by the source is acknowledged by the destination.
  - 3 **Take Down:** A four-way “Goodbye” handshake.



# Conversation Example

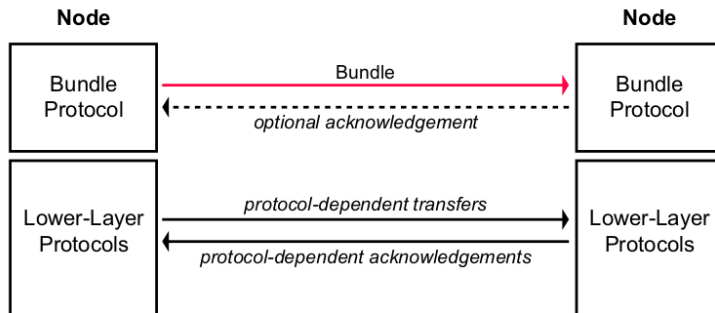


# Conversation Example



## A Non Conversational Protocol

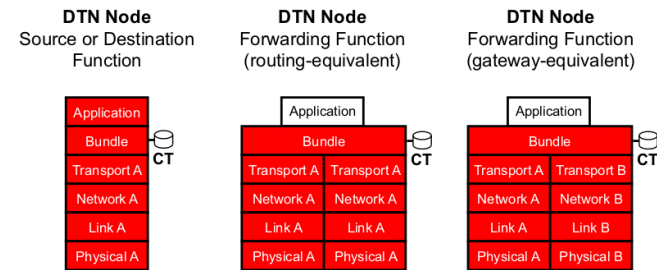
- DTN nodes communicate using simple sessions with minimal or no round-trips.
- Acknowledgements are optional.
- The Lower-layer protocols may be conversational (e.g., like TCP).



## DTN Node Roles

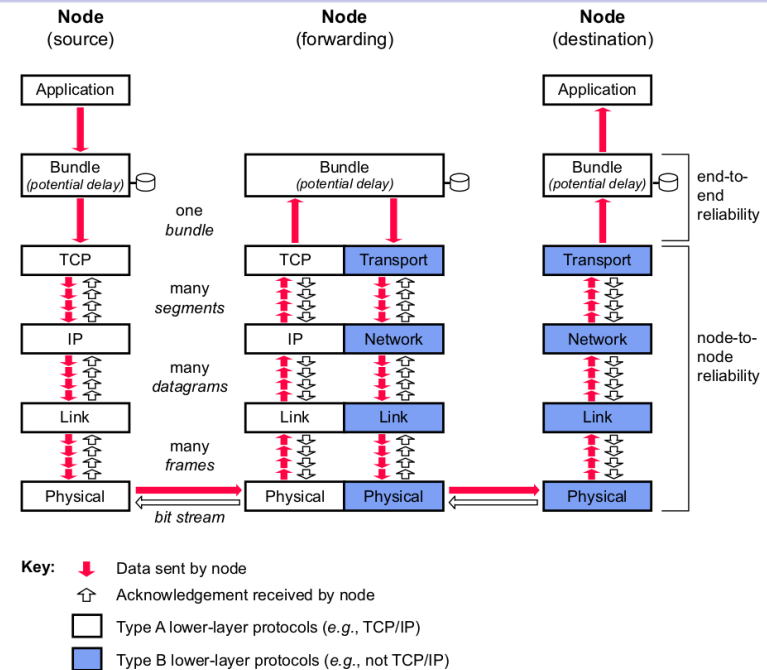
At any moment, a given node may act as a source, destination, or forwarder of bundles:

- **Source or Destination Function**
- **Forwarding Function**
  - 1 Routing-Equivalent Forwarding – implement same lower layer protocols.
  - 2 Gateway-Equivalent Forwarding – implement multiple stacks of lower layer protocols.



## Delay Isolation via Transport-Protocol Termination

- TCP protocol provides end-to-end (source-to-destination) reliability by retransmitting any segment that is not acknowledged
- The network, link, and physical protocols provide other types of data-integrity services.
- The bundle protocol relies on these lower-layer protocols to insure the reliability of communication.
- However, all DTN nodes terminate lower-layer transport protocols.
- **Problem:** The bundle protocol agents thus act as surrogates for end-to-end sources and destinations.
- **Opportunity:** Conversational lower-layer protocols are isolated by the bundle protocol from long delays elsewhere in the end-to-end path.



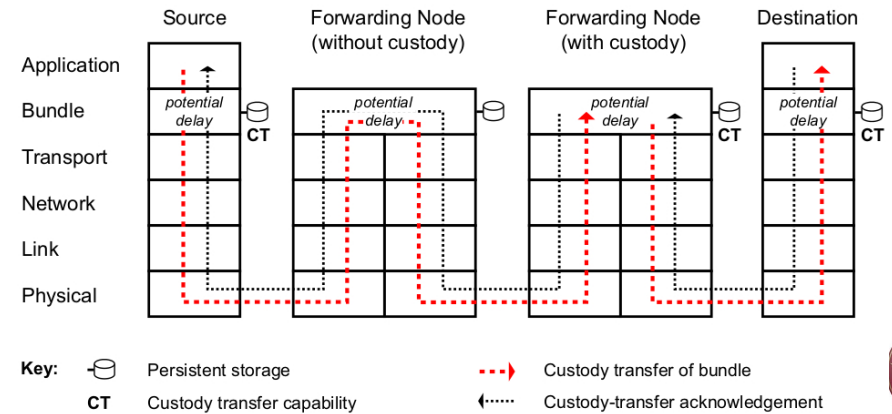
## Custody Transfers

- How to support node-to-node retransmission of lost or corrupt data ?
  - No single transport protocol typically operates end-to-end across a DTN.
  - End-to-end reliability can only be implemented at the bundle layer.
- Support node-to-node retransmission by means of **custody transfers**.
  - Custody transfers enhance end-to-end reliability,
  - but they do not guarantee it.
- Such transfers are arranged between successive nodes.
  - Not all successive nodes need to be custodian.
  - If the next successive node accepts custody, it returns an acknowledgment to the sender.



## Bundle Custodian

- A bundle custodian must store a bundle until either
  - 1 Another node accepts custody, or
  - 2 Expiration of the bundles time-to-live.

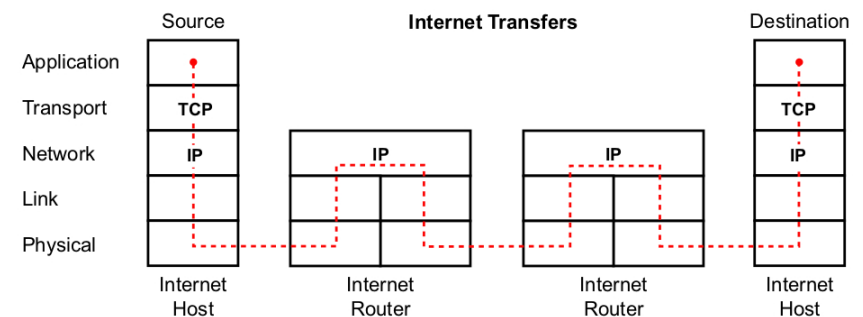


## Internet routing vs DTN routing

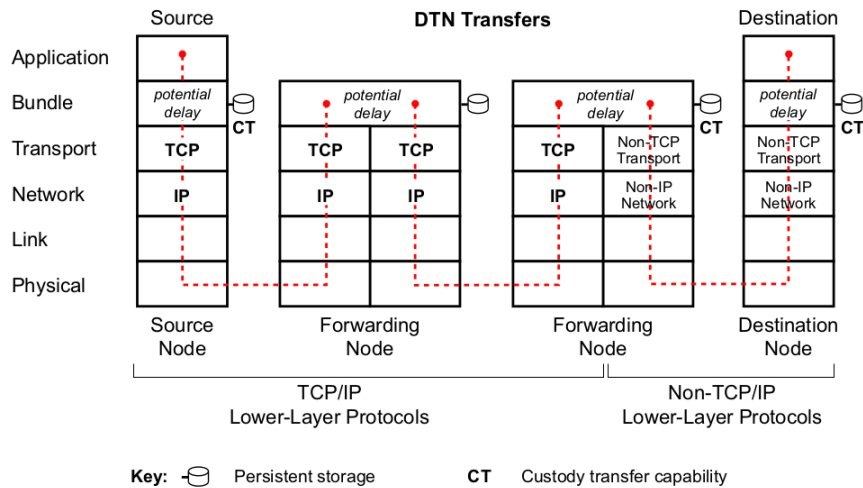
- On the Internet, the TCP and IP protocols are used throughout the network.
  - TCP operates at the end points of a path.
  - TCP manages reliable end-to-end delivery of TCP segments.
  - IP operates at all nodes on the path.
  - IP routes IP datagrams.
- In a DTN, all nodes implement both the bundle protocol and a lower-layer protocols.
  - Nodes that forward bundles can implement either the same or different lower-layer protocols on either side of the forwarding.
  - Nodes functions are comparable to Internet routers or gateways, respectively.



## Internet routing



## DTN routing



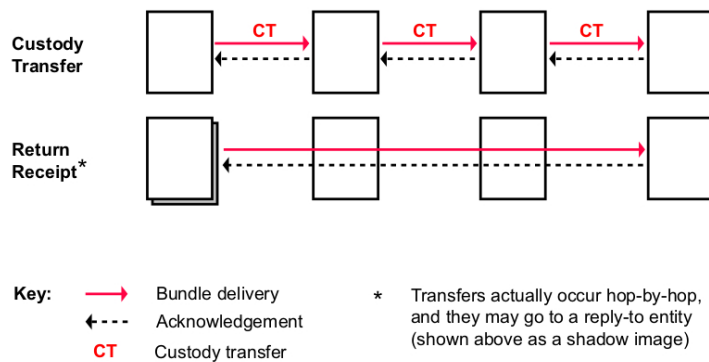
## Classes of Bundle Service

The bundle protocol provides six classes of service for a bundle:

- **Custody Transfer:** Delegation of retransmission responsibility by one node to another accepting node, so that the first node can recover its retransmission resources. The accepting node returns a custodial-acceptance acknowledgement to the previous custodian.
- **Return Receipt:** Confirmation by the destination to the source, or its reply-to entity, that the bundle has been received by the destination application. Reception by the source, or its reply-to entity, of the return receipt provides end-to-end assurance of delivery.
- **Priority of Delivery:** Bulk, Normal, or Expedited.
- **Time-to-Live**



## Classes of Bundle Service



## Endpoint IDs

- A bundle endpoint is a set of zero or more nodes that all identify themselves by the same endpoint ID.
- Common case: only one node has a given endpoint ID – called a **singleton endpoint**.
- Source nodes are always singleton endpoints or null (anonymous source) endpoints.
- Destination nodes may or may not be singleton endpoints.
- Endpoints may also be multicast (multiple destination nodes with the same endpoint ID) or null (no nodes).
- Endpoints may contain multiple nodes.
- Nodes may be members of multiple endpoints.



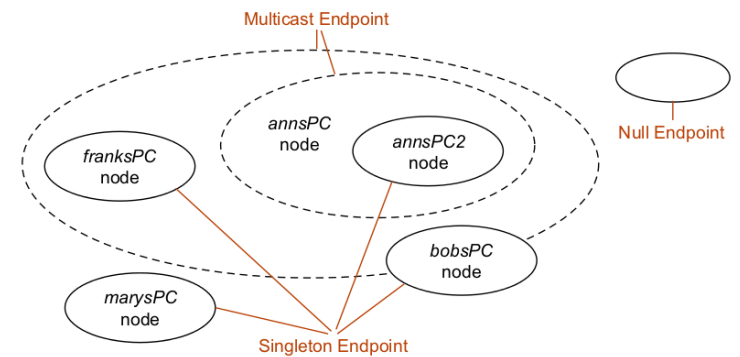
## Endpoint IDs

- An endpoint ID is a uniform resource identifier (URI) text string using the syntax:  
`<scheme_name>:<scheme-specific_part>`
- the scheme name is either `dtm` or `ipn`.
- The scheme-specific part comes in two flavors:
  - 1 Application-specific, used to identify a source or destination node, or
  - 2 Administrative, used when forwarding bundles from node to node.

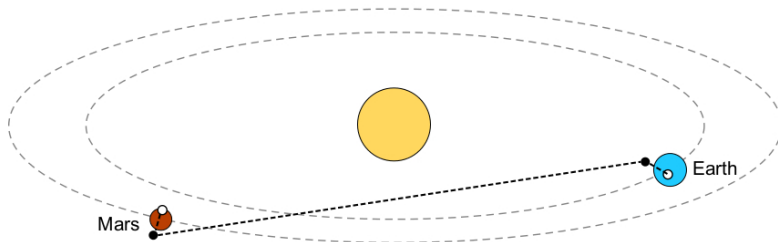
- `dtm://bobsPC/files` (application-specific)
- `dtm://bobsPC/` (administrative)
- `ipn:81.2` (application-specific)
- `ipn:81.0` (administrative)



## Naming Example



## A Simple Example



**Key:** - - - - - IPN link      ○ Source or destination node  
 ● Forwarding node



## A Simple Example

Node	Endpoint IDs	
Earth Source	<code>ipn:81.2</code>	(application-specific ID)
Earth Forwarding	<code>ipn:81.0</code>	(administrative ID)
	<code>ipn:49.0</code>	(administrative ID)
Mars Forwarding	<code>ipn:49.0</code>	(administrative ID)
	<code>ipn:65.0</code>	(administrative ID)
Mars Destination	<code>ipn:65.7</code>	(application-specific ID)



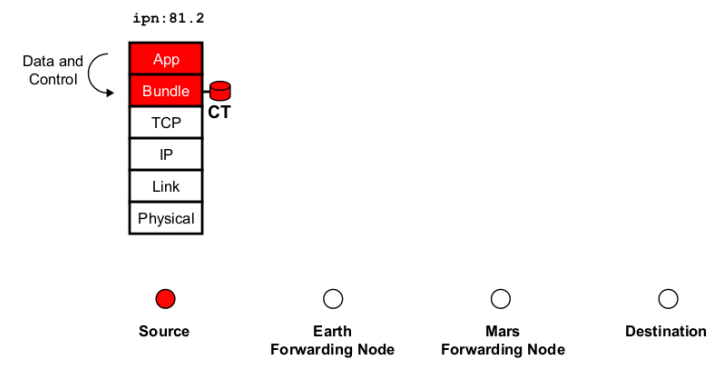


### A Simple Example

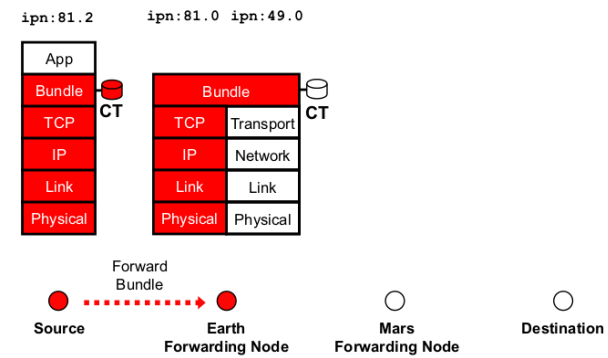
Item	Value
Source	ipn:81.2
Destination	ipn:65.7
Class of service	<ul style="list-style-type: none"> <li>• Custody transfer</li> <li>• Normal priority</li> <li>• Time-to-live = 36 hours</li> </ul>
User Data	Application-specific data, including instructions to the destination application for processing, storage, disposal, and error-handling. User data is not visible to bundle-protocol agents.



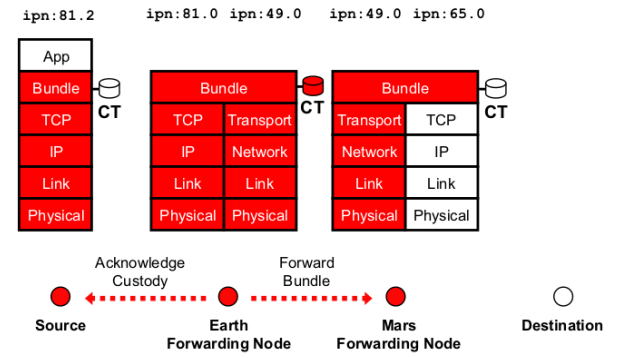
### A Simple Example



### A Simple Example

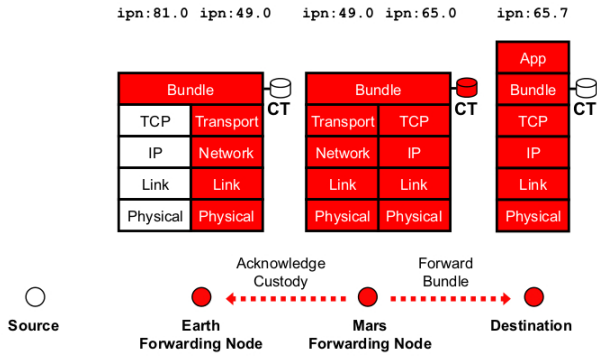


### A Simple Example



Example

# A Simple Example



Example

# A Simple Example

